

HOMEWORK 4

1. Define a field and distinguish between an infinite field and a finite field.
2. Use the extended Euclidean algorithm to find the inverse of $(x^4 + x^3 + 1)$ in $\text{GF}(2^5)$ using the modulus $(x^5 + x^2 + 1)$.
3. Create a table for addition and multiplication for $\text{GF}(2^4)$, using $(x^4 + x^3 + 1)$ as the modulus.
4. Using Table 1, perform the following operations:
 - a) $(100) \div (010)$
 - b) $(100) \div (000)$
 - c) $(101) \div (011)$
 - d) $(000) \div (111)$
5. Show how to multiply $(x^3 + x^2 + x + 1)$ by $(x^2 + 1)$ in $\text{GF}(2^4)$ using the algorithm in Table 2. Use $(x^4 + x^3 + 1)$ as modulus.
6. Show how to multiply (10101) by (10000) in $\text{GF}(2^5)$ using the algorithm in Table 3. Use $(x^5 + x^2 + 1)$ as modulus.

Example:

Find the result of multiplying $\mathbf{P_1} = (x^5 + x^2 + x)$, $\mathbf{P_2} = (x^7 + x^4 + x^3 + x^2 + x)$ in $\mathbf{GF}(2^8)$ with irreducible polynomial $(x^8 + x^4 + x^3 + x + 1)$.

Table 1: *Multiplication table for $\mathbf{GF}(2^3)$ with irreducible poynomial $(x^3 + x^2 + 1)$*

$\otimes$	000 (0)	001 (1)	010 (x)	011 (x+1)	100 (x ²)	101 (x ² +1)	110 (x ² +x)	111 (x ² +x+1)
000 (0)	000 (0)	000 (0)	000 (0)	000 (0)	000 (0)	000 (0)	000 (0)	000 (0)
001 (1)	000 (0)	001 (1)	010 (x)	011 (x+1)	100 (x ²)	101 (x ² +1)	110 (x ² +x)	111 (x ² +x+1)
010 (x)	000 (0)	010 (x)	100 (x ²)	110 (x ² +x)	101 (x ² +1)	111 (x ² +x+1)	001 (1)	011 (x+1)
011 (x+1)	000 (0)	011 (x+1)	110 (x ² +x)	101 (x ² +1)	001 (1)	010 (x)	111 (x ² +x+1)	100 (x ²)
100 (x ²)	000 (0)	100 (x ²)	101 (x ² +1)	001 (1)	111 (x ² +x+1)	011 (x+1)	010 (x)	110 (x ² +x)
101 (x ² +1)	000 (0)	101 (x ² +1)	111 (x ² +x+1)	010 (x)	011 (x+1)	110 (x ² +x)	100 (x ²)	001 (1)
110 (x ² +x)	000 (0)	110 (x ² +x)	001 (1)	111 (x ² +x+1)	010 (x)	100 (x ²)	011 (x+1)	101 (x ² +1)
111 (x ² +x+1)	000 (0)	111 (x ² +x+1)	011 (x+1)	100 (x ²)	110 (x ² +x)	001 (1)	101 (x ² +1)	010 (x)

 Table 2: *An efficient algorithm for multiplication using poynomials(Example)*

<i>Posers</i>	<i>Operation</i>	<i>New Result</i>	<i>Reduction</i>
$x^0 \otimes \mathbf{P}_2$		$x^7 + x^4 + x^3 + x^2 + x$	No
$x^1 \otimes \mathbf{P}_2$	$x \otimes (x^7 + x^4 + x^3 + x^2 + 1)$	$x^5 + x^2 + x + 1$	Yes
$x^2 \otimes \mathbf{P}_2$	$x \otimes (x^5 + x^2 + x + 1)$	$x^6 + x^3 + x^2 + x$	No
$x^3 \otimes \mathbf{P}_2$	$x \otimes (x^6 + x^3 + x^2 + x)$	$x^7 + x^4 + x^3 + x^2$	No
$x^4 \otimes \mathbf{P}_2$	$x \otimes (x^7 + x^4 + x^3 + x^2)$	$x^5 + x + 1$	Yes
$x^5 \otimes \mathbf{P}_2$	$x \otimes (x^5 + x + 1)$	$x^6 + x^2 + x$	No
$\mathbf{P}_1 \times \mathbf{P}_2 = (x^6 + x^2 + x) + (x^6 + x^3 + x^2 + x + 1) + (x^5 + x^2 + x + 1) = x^5 + x^3 + x^2 + x + 1$			

 Table 3: *An efficient algorithm for multiplication using n-bit words*

<i>Powers</i>	<i>Shift-Left Operation</i>	<i>Exclusive-Or</i>
$x^0 \otimes \mathbf{P}_2$		1001111
$x^1 \otimes \mathbf{P}_2$	00111100	$(00111100) \oplus (00011011) = \mathbf{00100111}$
$x^2 \otimes \mathbf{P}_2$	01001110	01001110
$x^3 \otimes \mathbf{P}_2$	10011100	10011100
$x^4 \otimes \mathbf{P}_2$	00111000	$(00111000) \oplus (00011011) = 00100011$
$x^5 \otimes \mathbf{P}_2$	01000110	01000110
$\mathbf{P}_1 \otimes \mathbf{P}_2 = (\mathbf{00100111}) \oplus (\mathbf{01001110}) \oplus (\mathbf{01000110}) = \mathbf{00101111}$		